## Finding Communities

Social Networks Analysis and Graph Algorithms
Prof. Carlos Castillo - https://chato.cl/teach

## Sources

- A. L. Barabási (2016). Network Science - Chapter 09
- D. Easly and J. Kleinberg (2010). Networks, Crowds, and Markets - Chapter 03
- F. Menczer, S. Fortunato, C. A. Davis (2020). A First Course in Network Science - Chapter 06
- URLs cited in the footer of slides


## Example with clear community structure



# Characterizing one community 

## Communities are connected and dense



Given a community $C$

Internal degree $k^{i n t}(C)$ considers only nodes inside the community

External degree $\boldsymbol{k}^{\text {ext }}(C)$ considers only nodes outside the community

$$
k_{i}=k_{i}^{\mathrm{int}}(C)+k_{i}^{\mathrm{ext}}(C)
$$

## Strong community



A community C is strong if every node $\boldsymbol{i}$ within the community satisfies:

$$
k_{i}^{\text {int }}(C)>k_{i}^{\text {ext }}(C)
$$

- Is the community of green nodes (dark green and light green) a strong community?
- What is the difference between dark green and light green nodes?


## Weak community



## A community C is weak if on

 aggregate nodes satisfy:$$
\sum_{i \in C} k_{i}^{\mathrm{int}}(C)>\sum_{i \in C} k_{i}^{\mathrm{ext}}(C)
$$

- All communities satisfying the strong property satisfy the weak one


## Exercise

Is community A strong, weak, both?
Is community $B$ strong, weak, both?
A community $C$ is strong if, for all nodes $i$ within the community:

$$
k_{i}^{\mathrm{int}}(C)>k_{i}^{\mathrm{ext}}(C)
$$

A community C is weak if:

$$
\sum_{i \in C} k_{i}^{\mathrm{int}}(C)>\sum_{i \in C} k_{i}^{\mathrm{ext}}(C)
$$




# Finding two communities: <br> network bisection 

## A graph that is easy to bisect



## Graph bisection: finding a minimal "cut"



# Simple exercise 

Cut size under bisection

- What is the size of the white-red cut?
- If node 9 goes to the red component, what is the size of the white-red cut?



## Finding multiple communities: a divisive method

## Hierarchical graph partitioning

Until there are edges in the graph
Find an edge $e$ that bridges two communities
Remove edge $e$

## The Girvan-Newman algorithm

- Repeat:
- Compute edge betweenness
- Remove edge with larger betweenness

b.

c.

d.




## Example: Karate Club




## Quantifying multiple communities: modularity

## Measuring a partition in a graph

- Modularity (or one of its variants) is a popular method to determine how good a partition is on a graph
- It compares the observed number of internal links in each partition, against the expected number of internal links if those internal links had been placed at random


## Modularity of a partition

$$
Q=\frac{1}{L} \sum_{C}\left(L_{C}-\frac{k_{C}^{2}}{4 L}\right)
$$

- $L=$ number of links in the network
- $L_{C}=$ number of internal links in community $C$
- $k_{C}=$ sum of degree of nodes in $C$


## Modularity of a partition (cont.)

$$
Q=\frac{1}{L} \sum_{C}\left(L_{C}-\frac{k_{C}^{2}}{4 L}\right) \longrightarrow \begin{aligned}
& \text { Expression in parenthesis } \\
& \text { is the difference between } \\
& \text { observed and expected } \\
& \text { internal links in } \\
& \text { community } C
\end{aligned}
$$

- $L=$ number of links in the network
- $L_{C}=$ number of internal links in community $C$
- $k_{C}=$ sum of degree of nodes in $C$
- $k^{2}{ }_{c} / 4 L=$ expected number of internal links in community $C$


## Where does $\mathrm{k}_{\mathrm{c}}{ }^{2} / 4 \mathrm{~L}$ comes from?

- A link "stub" is a connection between a link and a node
- There are $2 L$ stubs in a network
- There are as many stubs as the sum of the degree of nodes



## Modularity formula explained <br> $Q=\frac{1}{L} \sum_{C}\left(L_{C}-\frac{k_{C}^{2}}{4 L}\right)$

- There are $L_{C}$ internal links in $C$
- Total number of stubs in nodes in $C$ is $k_{C}$
- Total number of stubs in the network is $2 L$
- Probability of chosing two stubs in C: $\left(k_{c} / 2 L\right)^{2}=k_{c}{ }^{2} / 4 L^{2}$
- The expected number of links joining two stubs in $C$ is $L\left(k_{C}{ }^{2} / 4 L^{2}\right)=$ $k_{c}{ }^{2} / 4 \mathrm{~L}$
- The observed number is $L_{c}$
- $Q$ has a range: $Q \in[-1,+1]$
a.

c. SINGLE COMMUNITY

b. SUBOPTIMAL PARTITION

d. NEGATIVE MODULARITY

http://networksciencebook.com/chapter/9\#modularity

$$
Q=\frac{1}{L} \sum_{C}\left(L_{C}-\frac{k_{C}^{2}}{4 L}\right)
$$

## Exercise

- What is the modularity of the partition $\{0,1,2\},\{3,4,5\}$ ?

- What is the modularity of the partition $\{0,1,2,3\},\{4,5\}$ ?

$$
Q=\frac{1}{L} \sum_{C}\left(L_{C}-\frac{k_{C}^{2}}{4 L}\right)
$$



Pin board: https://upfbarcelona.padlet.org/chato/lak2lnp9m3jc1naj

## Summary

## Things to remember

- Strong and weak community
- The concept of "cut" in graph bisection
- Girvan-Newman's algorithm
- Modularity


## Practice on your own

- Check the modularity computations in the example on the slide marked $\boldsymbol{*}$ : (a) optimal partitioning into two communities, (b) suboptimal partitioning into two communities, (c) all the nodes in a single community, (d) one community per node
- You can check your answers with networkx.algorithms.community.modularity

