

# Graph Theory: Basics

## Introduction to network Science

Instructor: Michele Starnini — <https://github.com/chatox/networks-science-course>



Universitat  
Pompeu Fabra  
*Barcelona*

# Network Science or Graph Theory?

	Field	When	What	How
<b>Graphs</b>	mathematics (computer science)	1960-70	Theory	Structural properties
<b>Networks</b>	physics	2000- present	Applications	Complex systems, Dynamical processes

Nobody cares! Completely equivalent

# What is (modern) Network Science

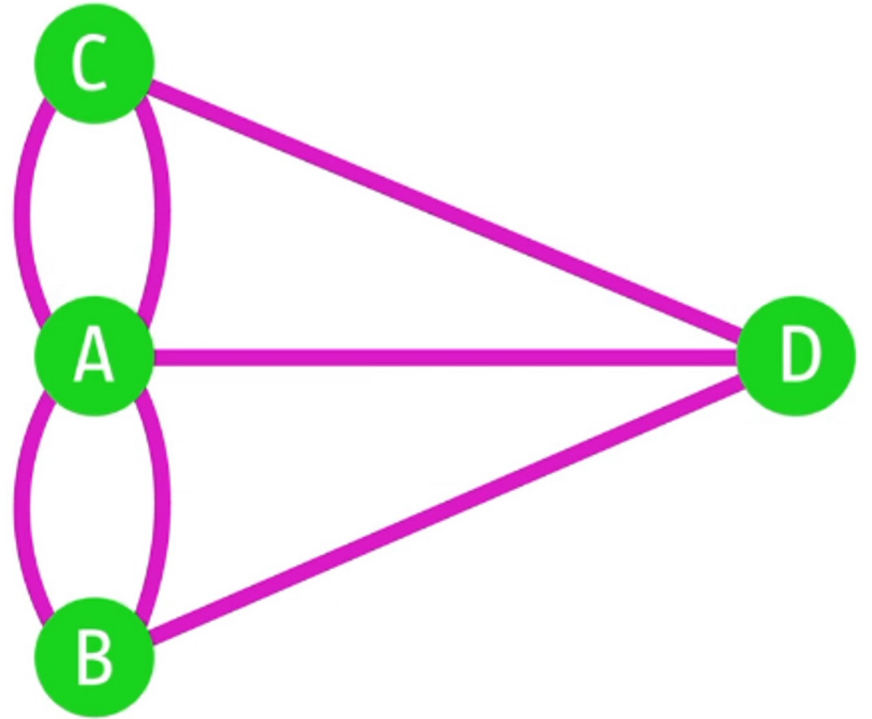
- **Mathematical formalism** from graph theory
- New insights from **statistical physics**
- Recent availability of very **large databases** on networked systems
- **Computational resources** to analyze them

# Contents

- Directed, weighted, ... many kind of graphs!
- Density & Sparsity
- Degrees
- Adjacency matrices

# Notation for a graph

- $G = (V, E)$ 
  - $V$ : nodes or vertices
  - $E$ : links or edges
- $|V| = N$  size of graph
- $|E| = L$  (or  $E$ ) number of links



# Subgraph

- Given  $G = (V, E)$
- A **subgraph** induced by a nodeset  $S$  is the graph  $G' = (S, F)$  defined by:
  - nodes in  $S$
  - edges in  $F = \{ (u, v) \in E \text{ s.t. } u \in S \text{ and } v \in S \}$

# Directed vs undirected graphs

- In an undirected graph
  - $E$  is a symmetric relation

$$(u, v) \in E \Rightarrow (v, u) \in E$$

- In a directed graph, also known as “digraph”
  - $E$  is not a symmetric relation

$$(u, v) \in E \not\Rightarrow (v, u) \in E$$

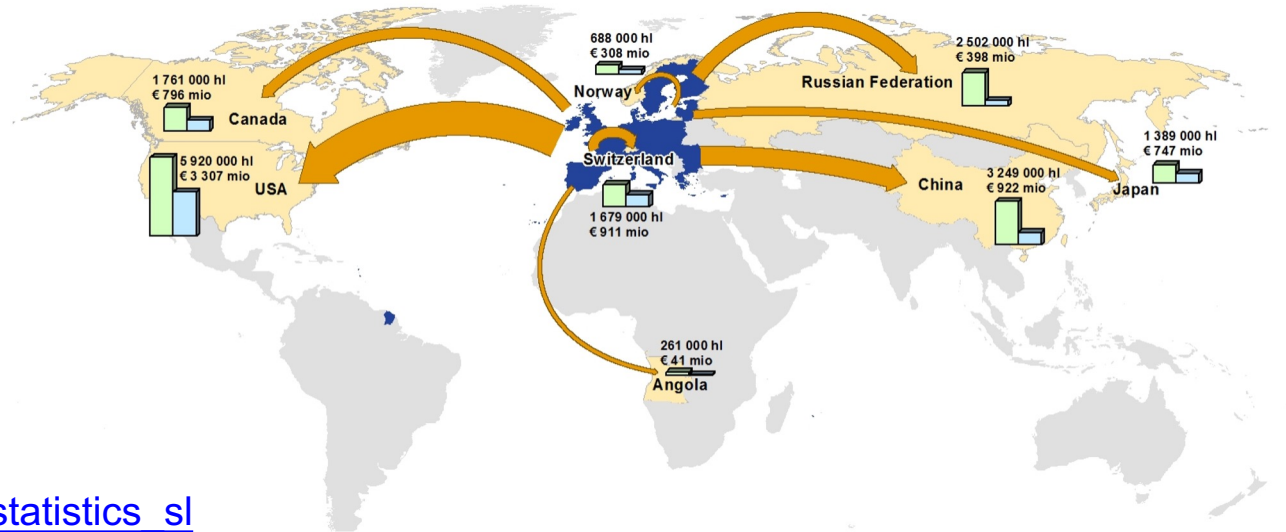
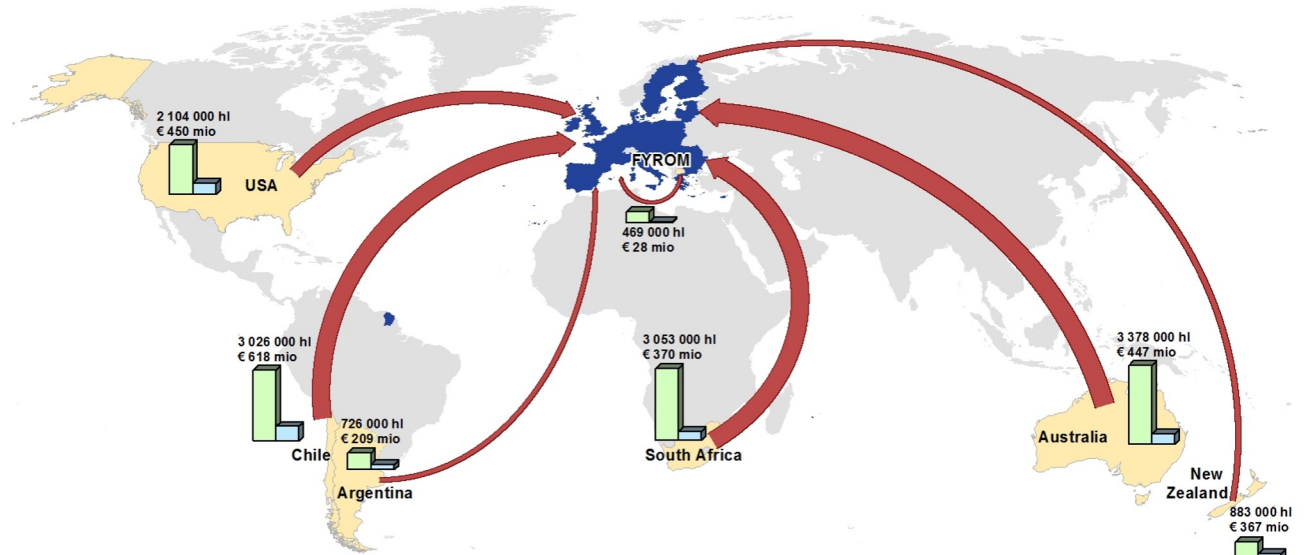
# Weighted vs unweighted graphs

- In a weighted graph edges have **weights** denoting the strength or importance of a connection
- When drawing, typically larger weights are drawn with ticker lines



# Weighted networks

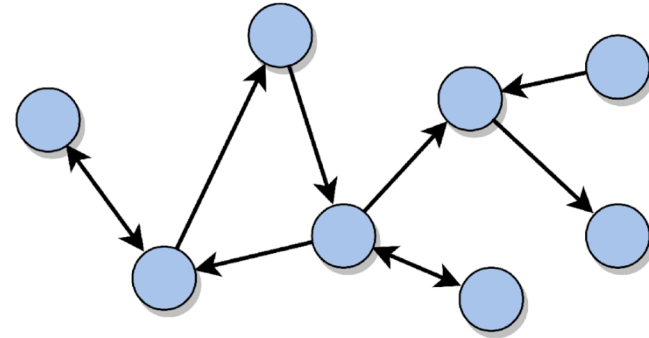
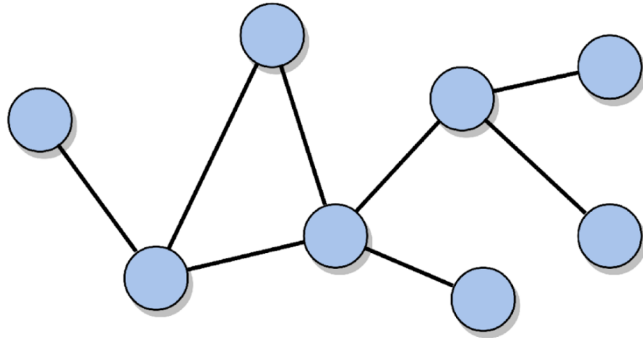
EU imports (top)  
and exports (bottom)  
of wine



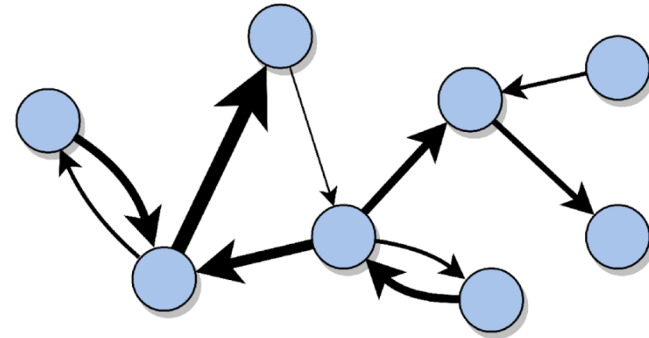
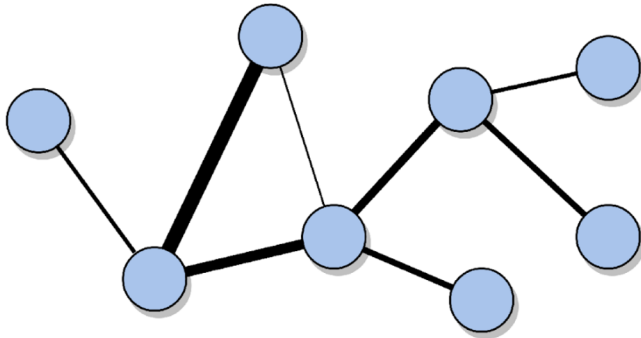
Undirected

Directed

Unweighted



Weighted



# Example graphs we will use

Network	N	E
Zachary's Karate Club (karate.gml)	34	78
Game of Thrones (got-relationships.csv)	84	216
US companies ownership	1351	6721
Marvel comics (hero-network.csv)	6K	167K

# Density and sparsity

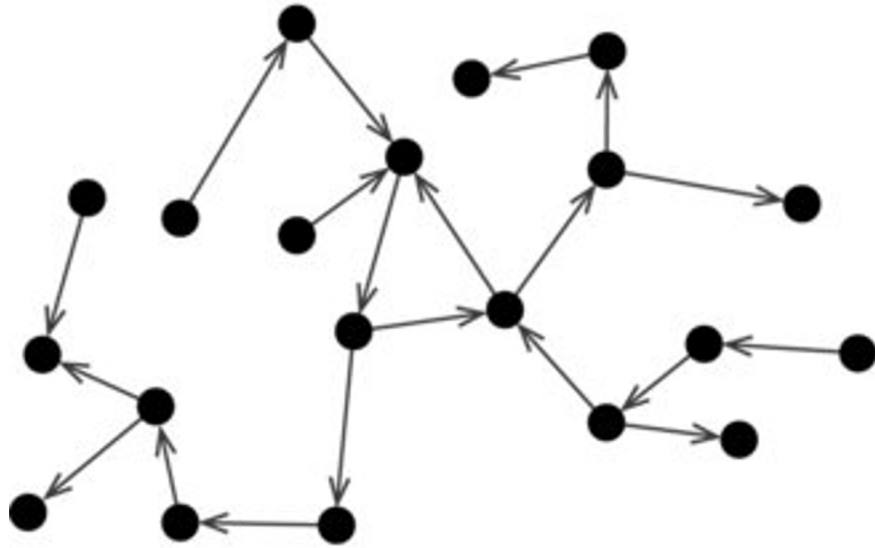
- Network size  $N$  = number of nodes
- $L$  = number of links
- Maximum possible number of links:

$$L_{max} = \binom{N}{2} = \frac{N(N-1)}{2}$$

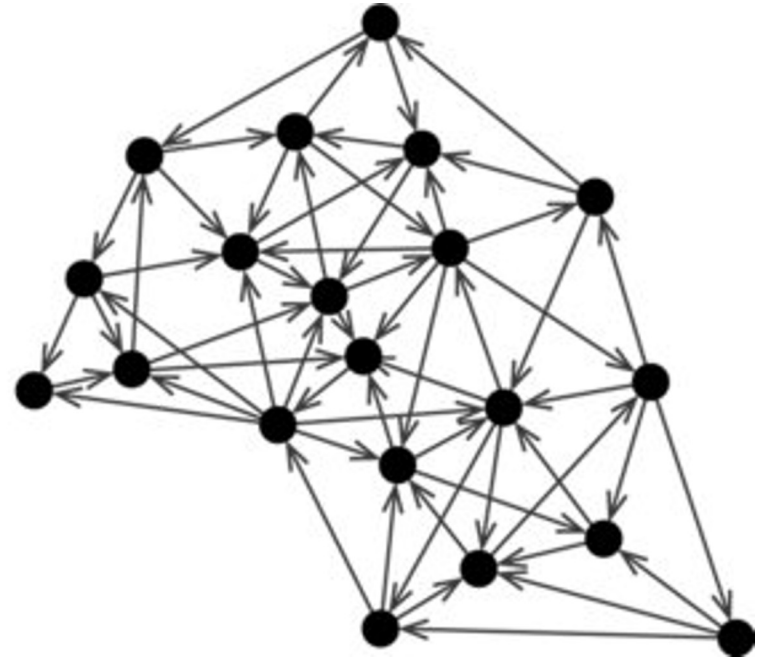
- Density:  $d = \frac{L}{L_{max}} = \frac{2L}{N(N-1)}$

- The network is **sparse** if  $d \ll 1$

# Sparse network



# Dense network



# Example: Facebook

Rough orders-of-magnitude approximations:

- $N \approx 10^9$
- $L \approx 10^{13} \times N$
- $d \approx L / N^2 \approx 10^{13} N / N^2 \approx 10^{13} / 10^9 = 10^4$
- Most (but not all) real-world networks are similarly sparse because the number of links scales proportionally to  $N$ , whereas the maximum scales with  $N^2$

# Complex networks are sparse

- Theoretically  $L_{\max} = \binom{N}{2} = \frac{N(N-1)}{2}$
- Most real networks are sparse, i.e.,  
 $L \ll L_{\max}$

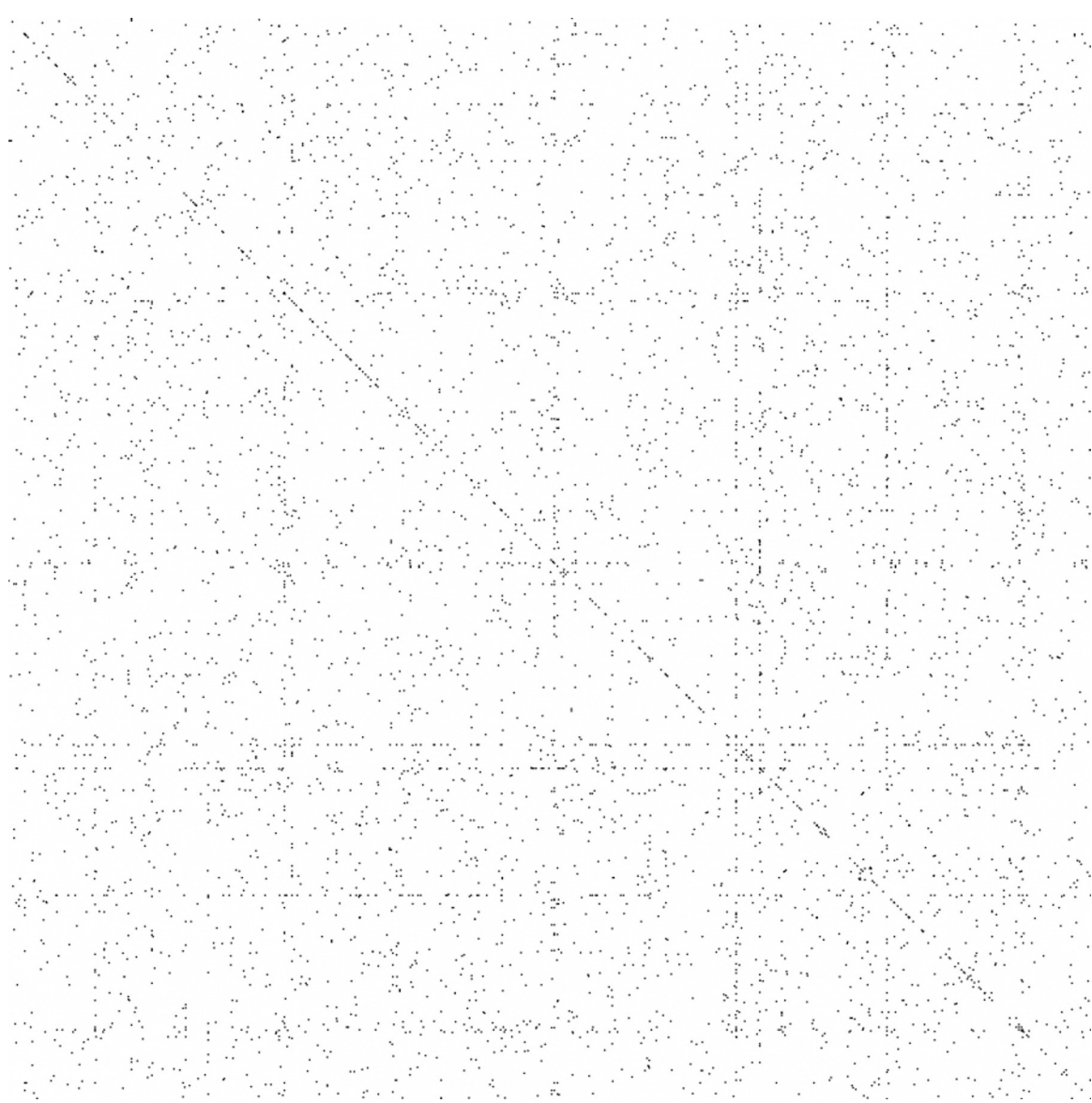
L is the number of links in the network, N is the number of nodes on it

# How sparse are some networks?

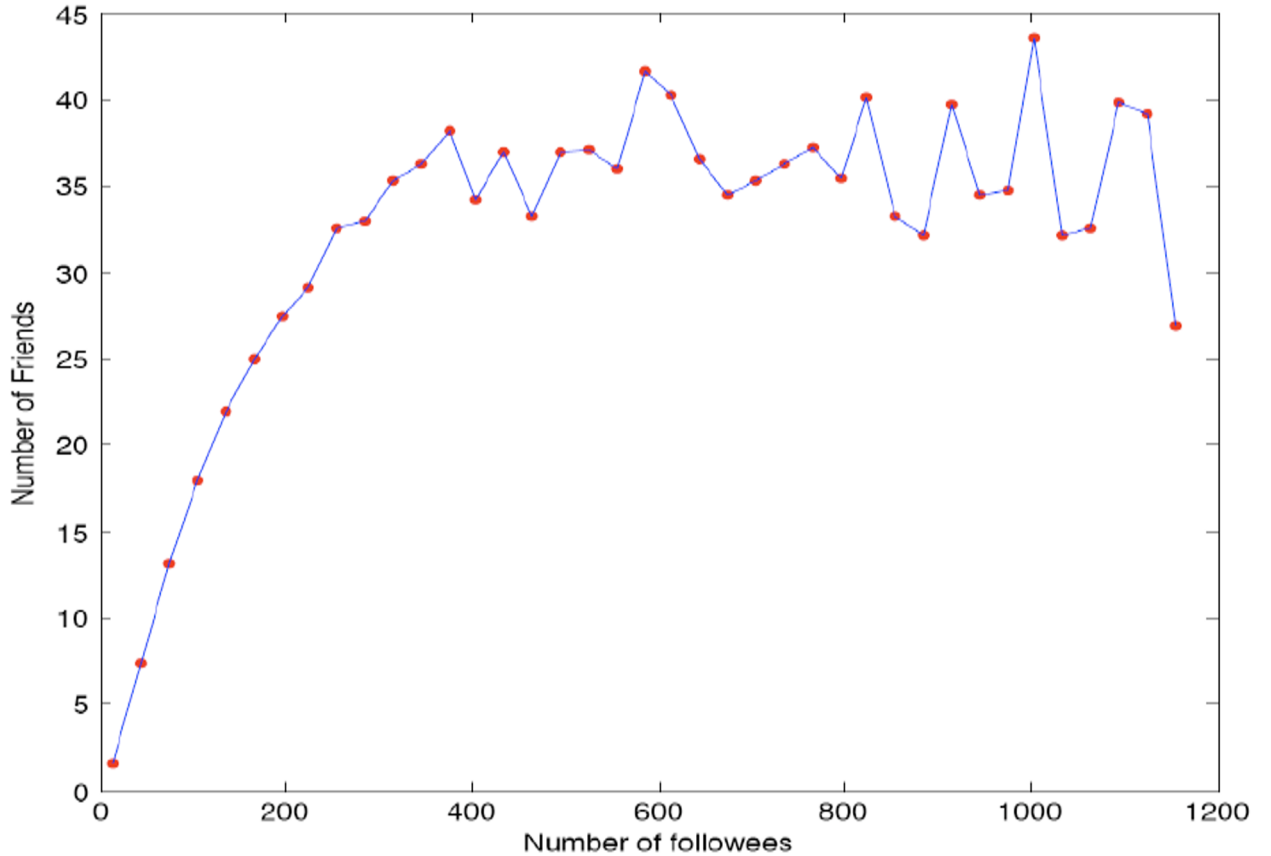
Network	$ V $	$ E $	Max $ E $
Zachary's Karate Club	34	78	561
Game of Thrones	84	216	3496
US companies ownership	1351	6721	911K
Marvel comics	6K	167K	17M



**Example:  
protein  
interaction  
network  
( $N=2K$ ,  $L=3K$ )**



Example: people you follow on Twitter (followees) vs people you have sent at least two messages to (“friends”)



# Why are networks sparse?

- Different mechanisms, think about it from the node perspective:
  - How many items **could** the node be connected to
  - Would it be **realistic** to connect to a large fraction of them?
- In social networks, Dunbar's number ( $\approx 150$ )

# DUNBAR'S NUMBER : 150

TYPICAL NUMBER OF PEOPLE WE CAN KEEP TRACK OF AND  
CONSIDER PART OF OUR ONGOING SOCIAL NETWORK

150  
TRIBE

50  
CLAN

15  
SUPER  
FAMILY

5  
CLOSE  
FRIENDS

← WEAKER TIES → MORE INVESTMENT IN RELATIONSHIP →

DANG! NOW,  
WHAT WAS THEIR  
NAME AGAIN?



# Degree

Node  $i$  has degree  $k_i$

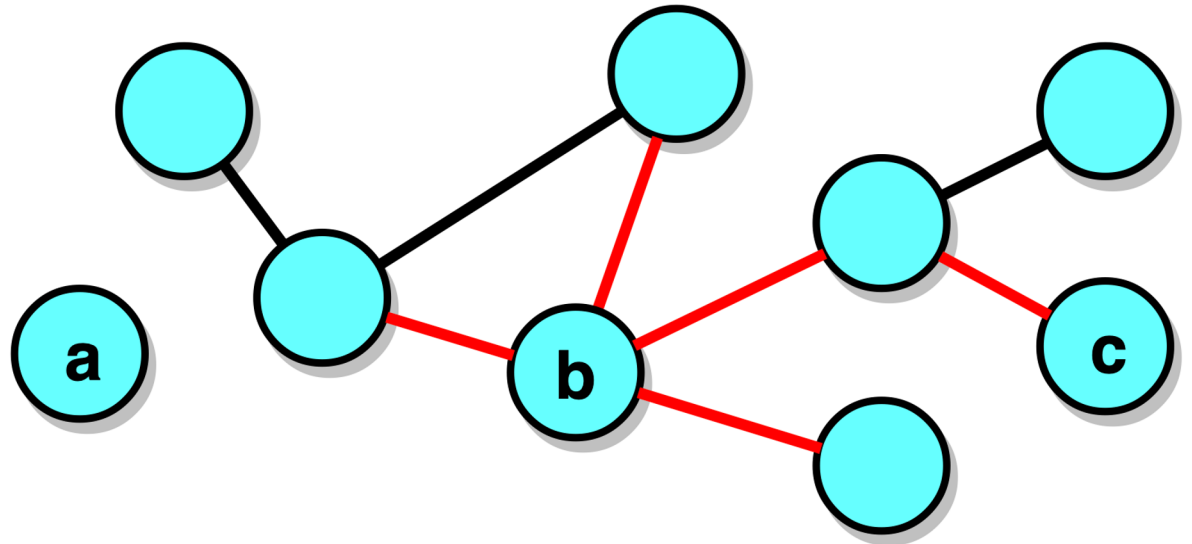
This is the number of links incident on this node

The total number of links  $L$  is given by

$$L = \frac{1}{2} \sum_{i=1}^N k_i$$

Average degree

$$\langle k \rangle = \frac{1}{N} \sum_{i=1}^N k_i = \frac{2L}{N}$$



# Average degree & density

$$d = \frac{\langle k \rangle}{N - 1} = \frac{\langle k \rangle}{k_{max}}$$

- Let us revisit the Facebook example
- Since  $N$  is very large,  $N - 1$  can be approximated by  $N$

$$d = \frac{\langle k \rangle}{N - 1} \approx \frac{\langle k \rangle}{N} \approx \frac{10^3}{10^9} = 10^{-6}$$

# In directed graphs

We distinguish **in-degree** from **out-degree**

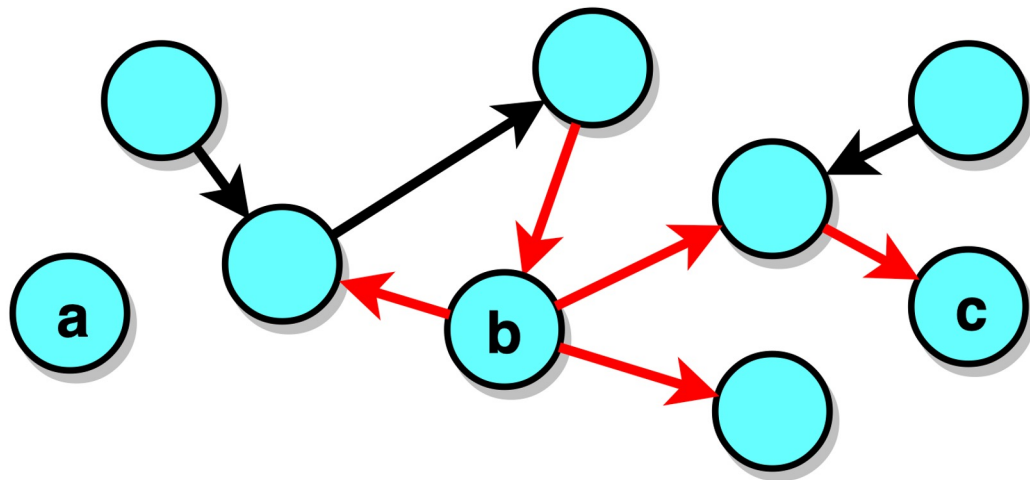
Incoming and outgoing links, respectively

Degree is the sum of both

$$k_i = k_i^{\text{in}} + k_i^{\text{out}}$$

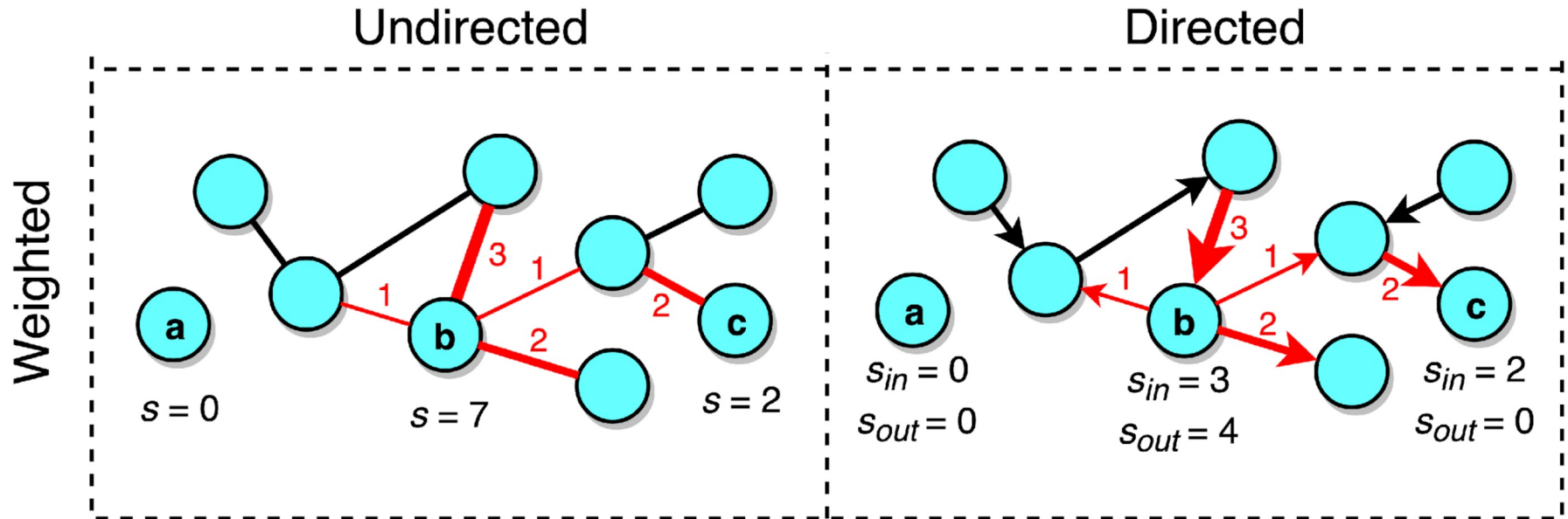
Counting total number of links:

$$L = \sum_{i=1}^N k_i^{\text{in}} = \sum_{i=1}^N k_i^{\text{out}}$$



# In weighted graphs

- We speak of “weighted degree” or “strength”



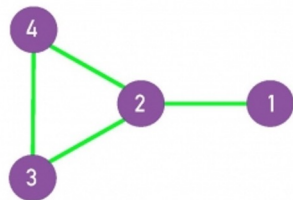


# Degree distribution

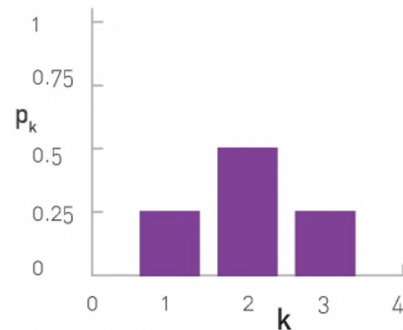
- If there are  $N_k$  nodes with degree  $k$
- The **degree distribution** is given by  $p_k = \frac{N_k}{N}$
- The average degree is then  $\langle k \rangle = \sum_{k=0}^{\infty} k p_k$

# Degree distribution; two toy graphs

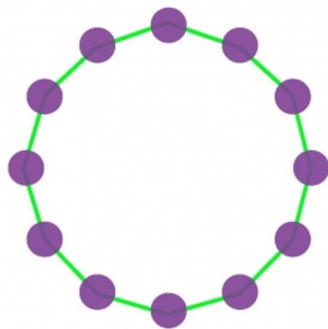
a.



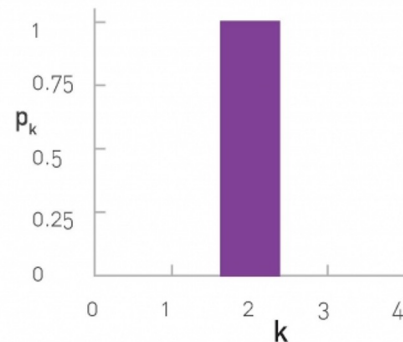
b.



c.

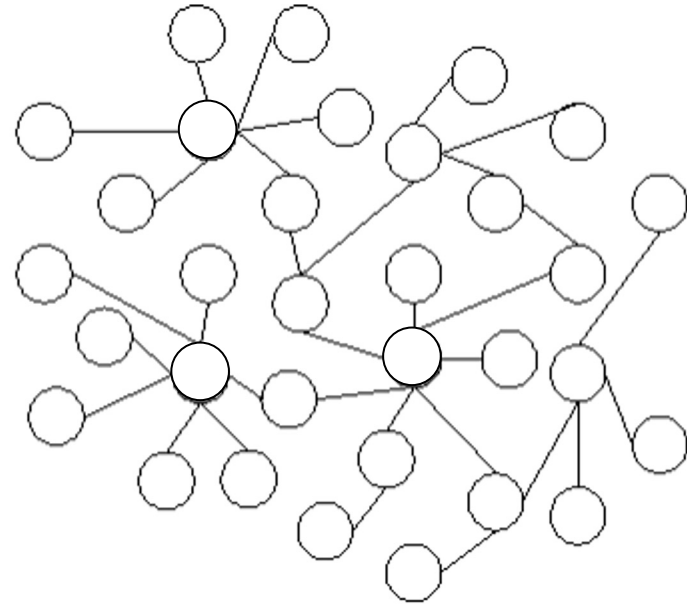
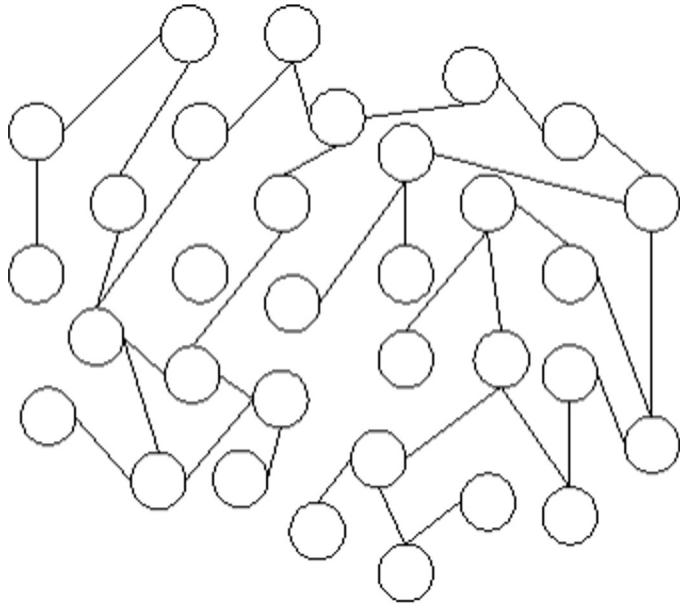


d.



# Exercise

Draw the degree distribution of these graphs

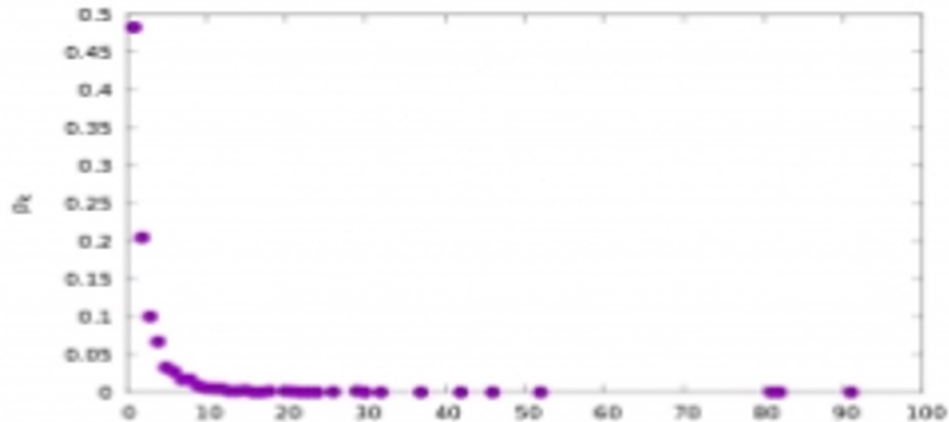
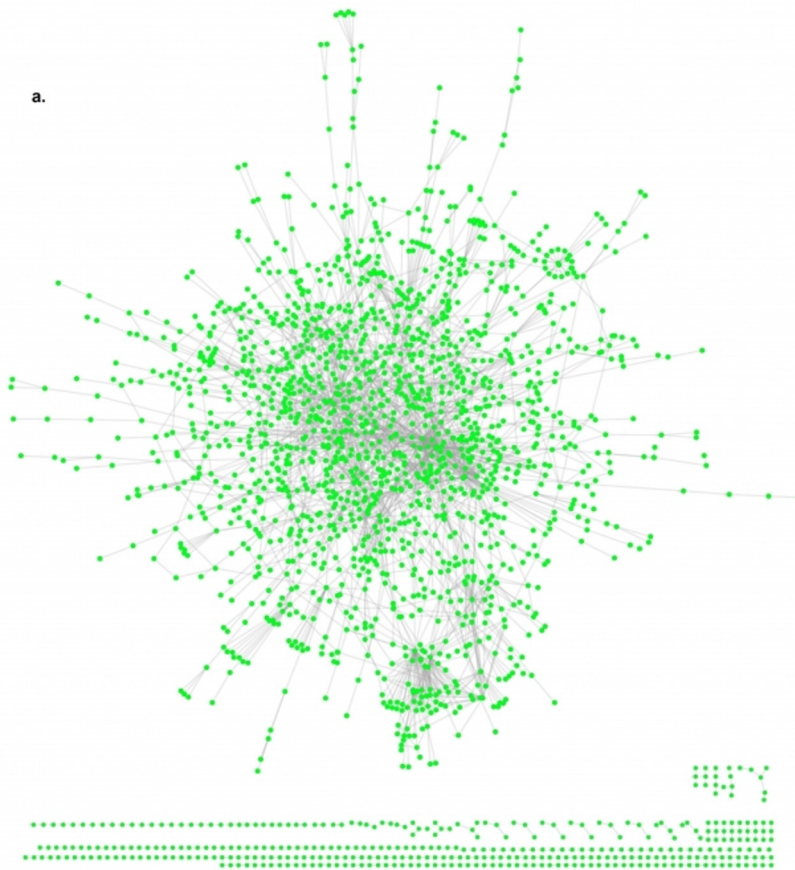


Spreadsheet links: <https://upfbarcelona.padlet.org/chato/shyq9m6f2g2dh1bw>



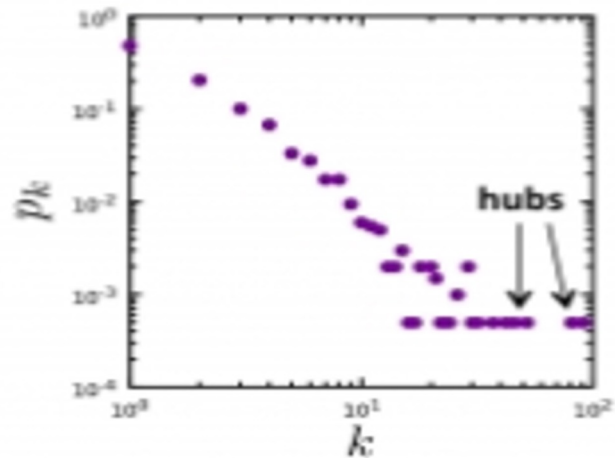
# Real graphs

a.



Linear  
scale

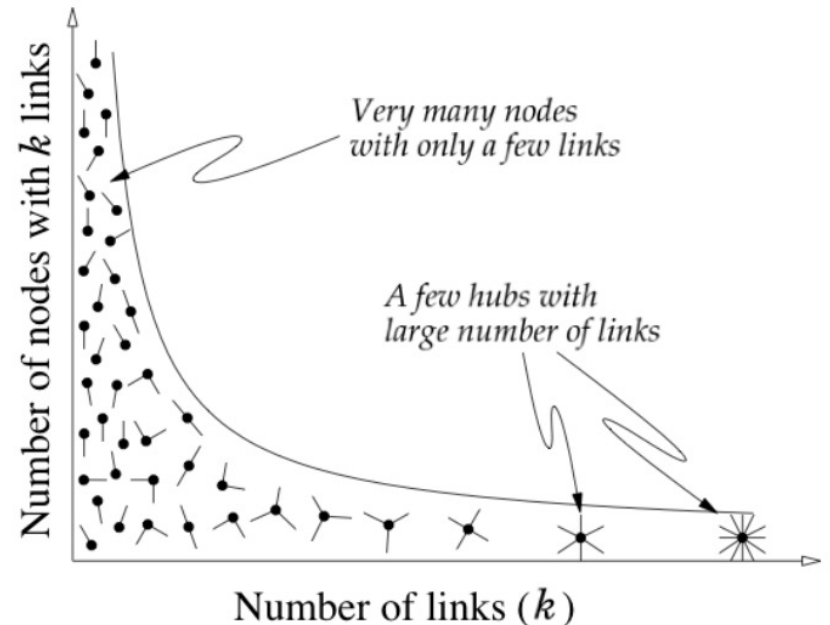
c.



Log-log  
scale

# Degree distribution $P(k)$

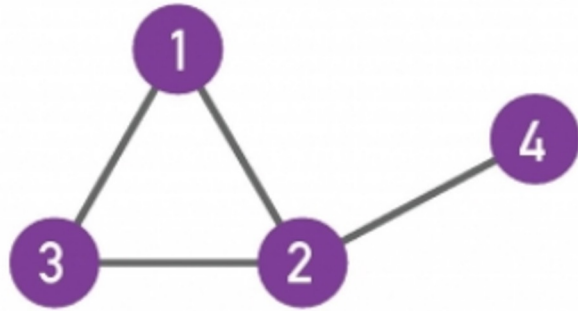
- Most important statistical property of networks
- Nodes with large degree = hubs, nodes with small degree = leaves
- When hubs are present, the  $P(k)$  is heterogeneous
- Most real (complex) networks are heterogeneous



# Adjacency matrix

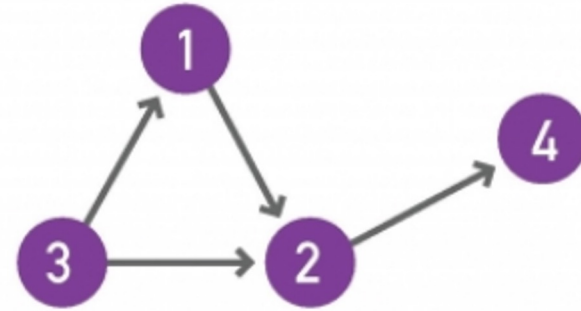
- A is the **adjacency matrix** of  $G = (V, E)$  iff:
  - A has  $|V|$  rows and  $|V|$  columns
  - $A_{ij} = 1$  if  $(i,j) \in E$
  - $A_{ij} = 0$  if  $(i,j) \notin E$
- **$A_{ij}$  always means row i, column j**
  - Sometimes Barabási's book has this wrong

# Examples



Undirected graph

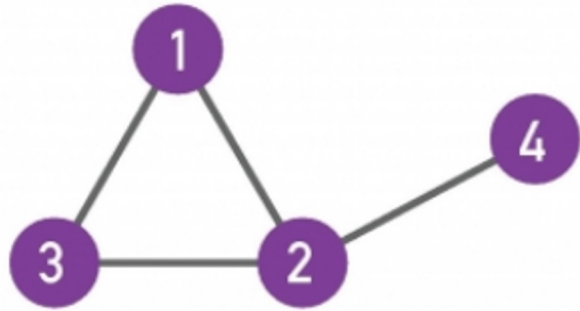
$$A_{ij} = \begin{matrix} & 1 & 2 & 3 & 4 \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & 0 & 1 & 1 & 0 \\ & 1 & 0 & 1 & 1 \\ & 1 & 1 & 0 & 0 \\ & 0 & 1 & 0 & 0 \end{matrix}$$



Directed graph

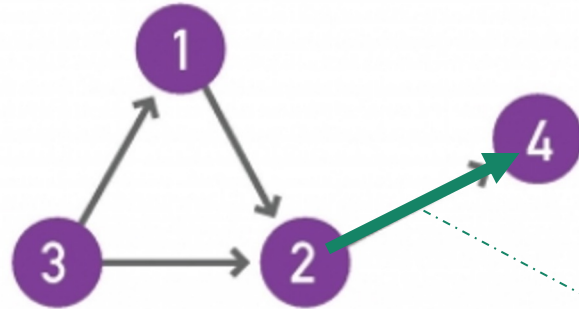
$$A_{ij} = \begin{matrix} & 1 & 2 & 3 & 4 \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & 0 & 1 & 0 & 0 \\ & 0 & 0 & 0 & 1 \\ & 1 & 1 & 0 & 0 \\ & 0 & 0 & 0 & 0 \end{matrix}$$

# $A_{ij}$ always means row $i$ , column $j$



Undirected graph

$$A_{ij} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{matrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{matrix} \end{matrix}$$



Directed graph

$$A_{ij} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{matrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{matrix} \end{matrix}$$

Row 2  
Column 4



# Properties of adjacency matrices

- $G$  is undirected  $\Leftrightarrow A$  is symmetric
- $G$  has a **self-loop**  
 $\Leftrightarrow A$  has a non-zero element in the diagonal
- $G$  is **complete**  $\Leftrightarrow A_{ij} \neq 0$  (except if  $i=j$ )

# Quick Exercise

- In terms of  $A$ , what is the expression for:

$$k_i^{\text{in}} =$$

$$k_i^{\text{out}} =$$

# Weighted networks

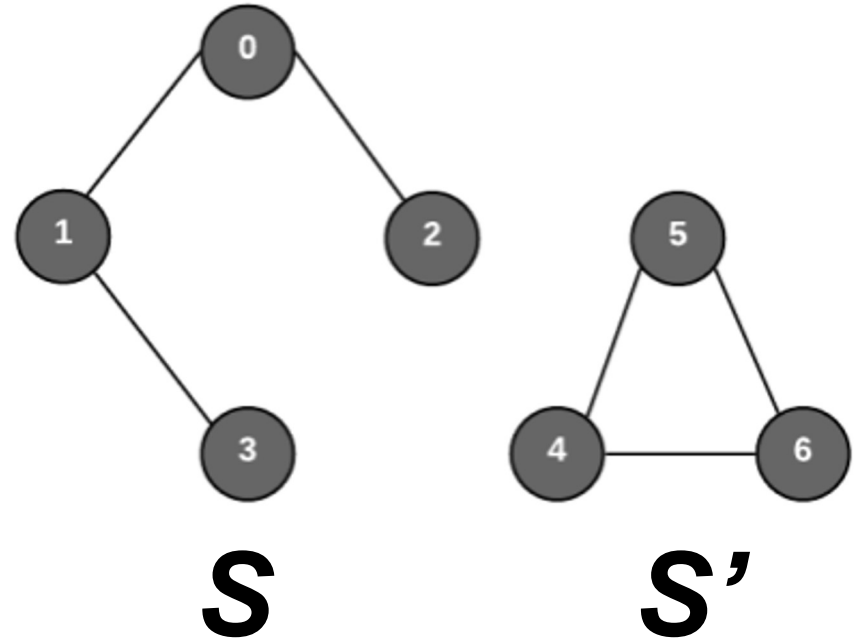
- Element  $w_{ij}$  indicates the **weight of the link** from node  $i$  to node  $j$
- The link from node B to node C has a weight of 2
- Out-strength (weighted out-degree):  $s_i^{out} = \sum_j A_{ij}$
- In-strength (weighted in-degree):  $s_i^{in} = \sum_j A_{ji}$
- Total weight:  $W = \sum_i s_i^{in} = \sum_i s_i^{out} = \sum_{ij} A_{ij}$

$$\begin{matrix} & A & B & C & D & E & F \\ \begin{matrix} A \\ B \\ C \\ D \\ E \\ F \end{matrix} & \begin{pmatrix} 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 2 & 1 & 3 & 1 & 1 & 0 \end{pmatrix} \end{matrix}.$$

# If a graph is disconnected

Disconnected graphs  
have adjacency  
matrices with **block  
structure**

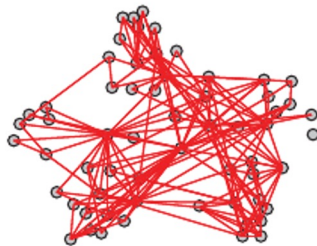
$$A = \begin{bmatrix} S & 0 \\ 0 & S' \end{bmatrix}$$



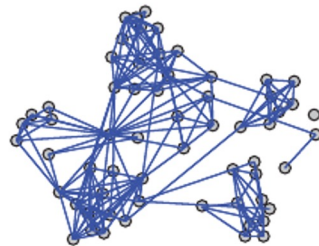
# Beyond simple graphs

# Some networks are multi-layer

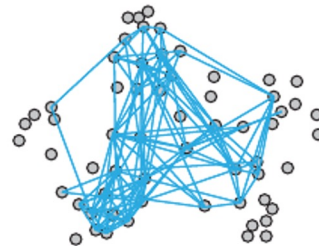
- Multi-layer graphs have different edges **over the same nodes**
- Example graph of 61 employees of a university department, with links indicating: coworking, having lunch together, being connected on Facebook, being actual friends, or being co-authors



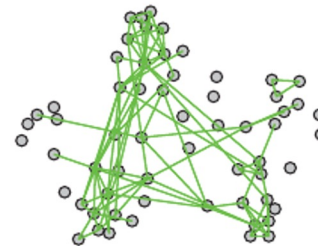
(a) Work.



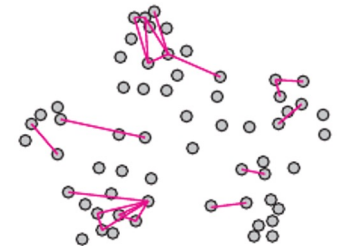
(b) Lunch.



(c) Facebook.



(d) Friend.



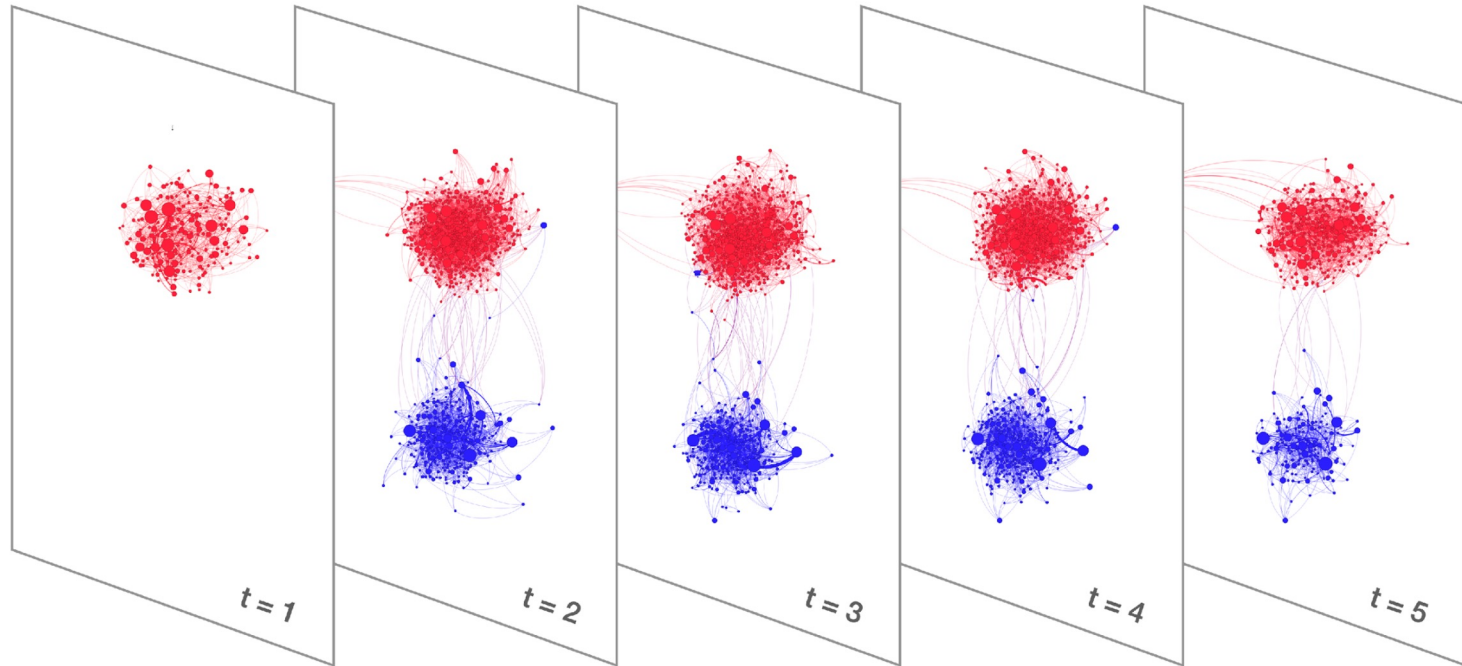
(e) Coauthor.

# Some networks are time-evolving

**Temporal, or  
“time-evolving”**

networks

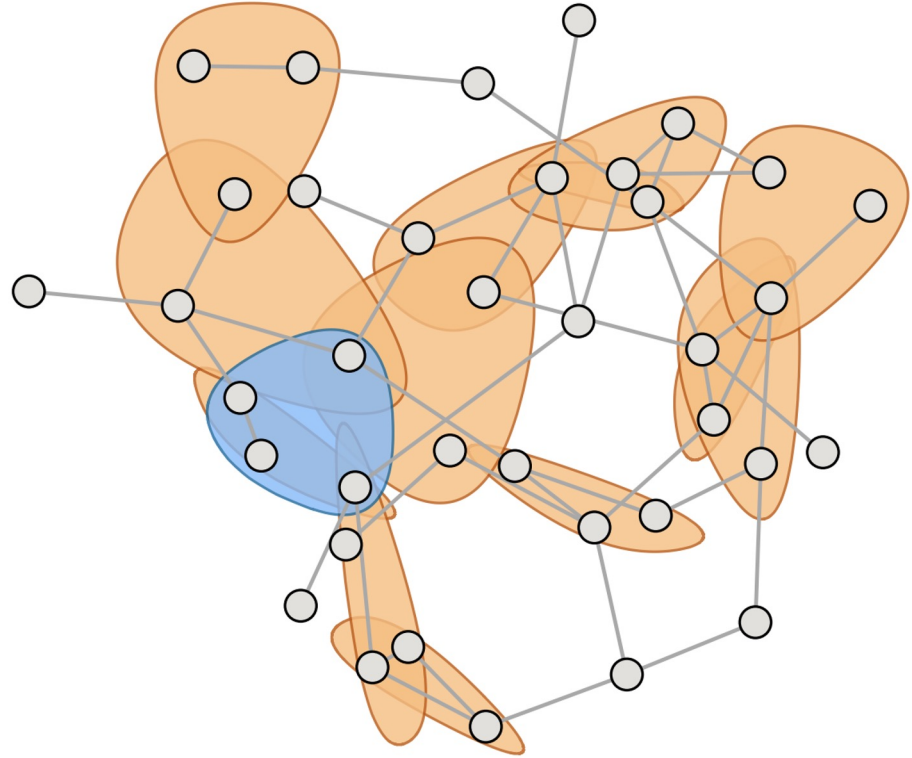
At each timestep  
there are new  
nodes and/or  
edges (and/or  
deletions)



# Some networks are higher-order

## Higher-order networks, or “hypergraphs”:

Hyper-links involve more than two nodes

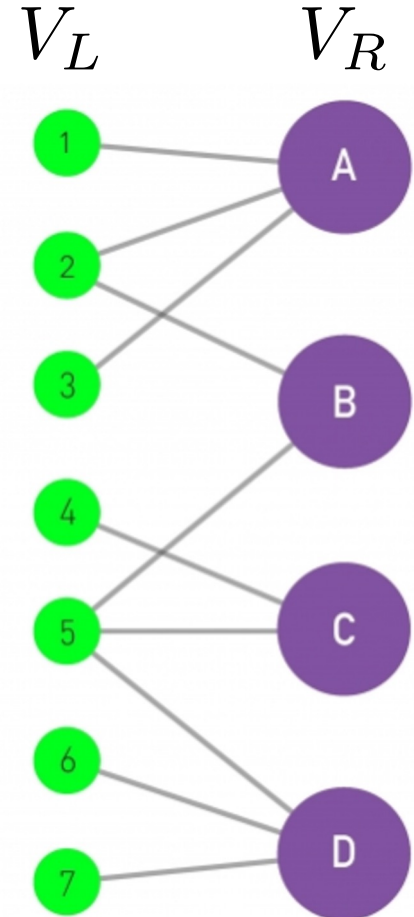




# Some networks are **bi-partite**

- A **bipartite** graph is a graph
- $G = (V, E)$  such that

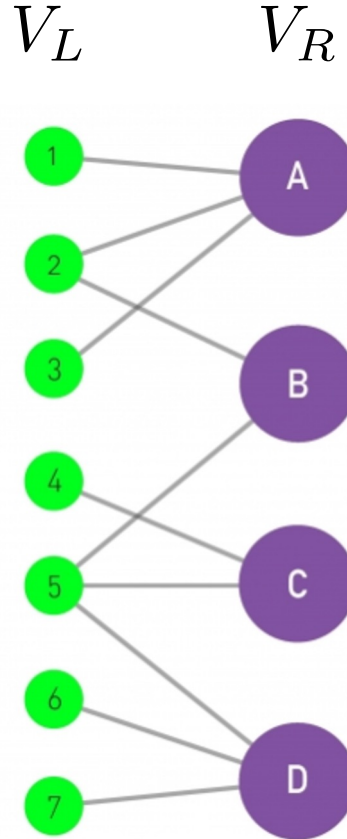
$$V = V_L \cup V_R, V_L \cap V_R = \emptyset, E \subseteq V_L \times V_R$$



# Exercise: project a bipartite network

?

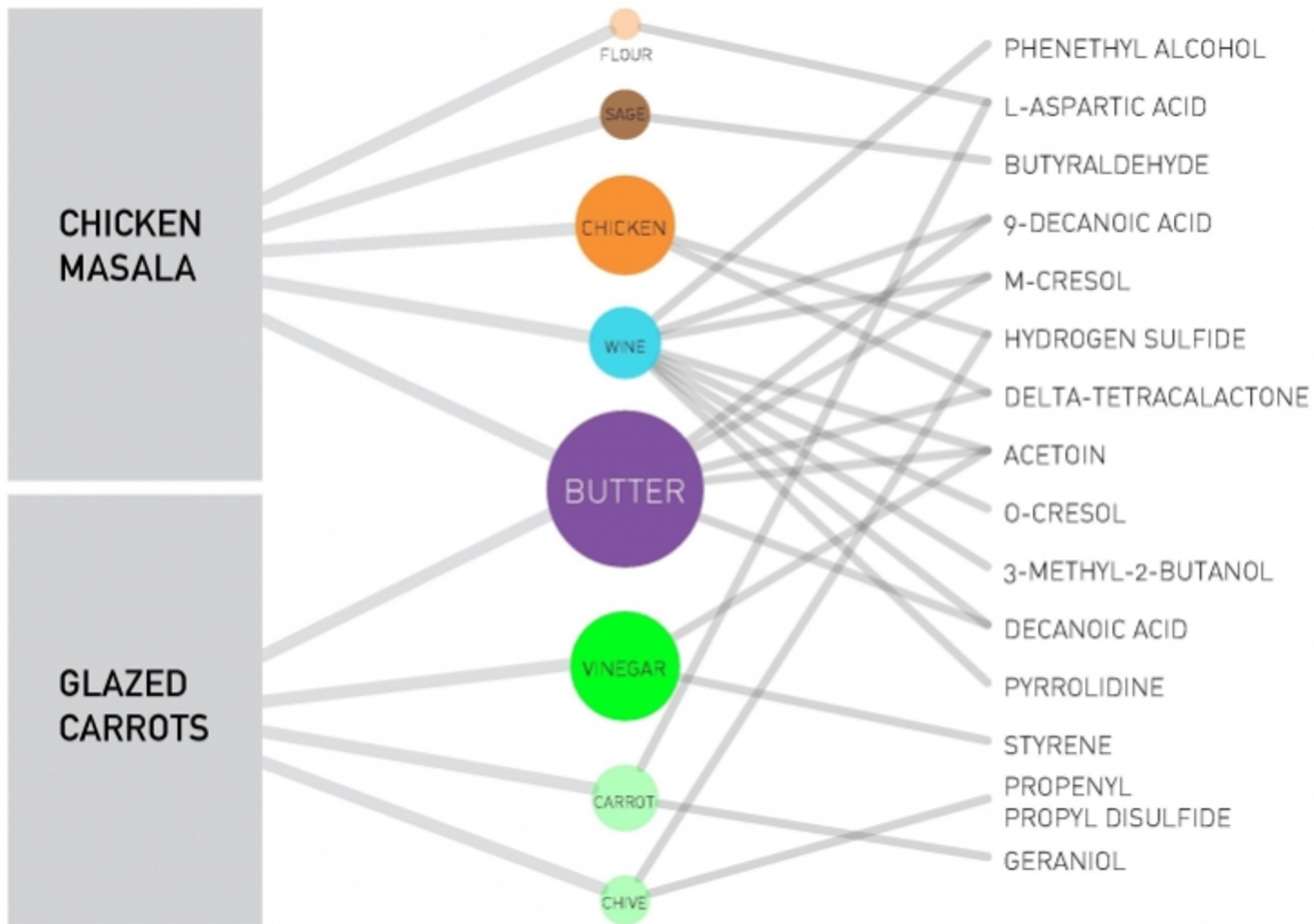
Left projection:  
graph where  
nodes  
are 1, 2, ..., 7  
and  
nodes are  
connected  
if they share a  
neighbor



?

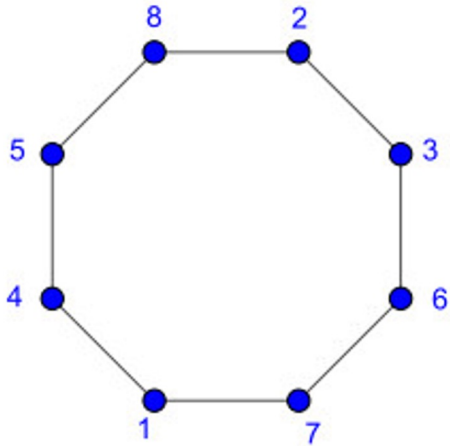
Right  
projection:  
graph where  
nodes  
are A, B, ..., D  
and  
nodes are  
connected  
if they share a  
neighbor

# Tripartite network

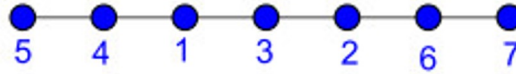


# Some graphs have a name

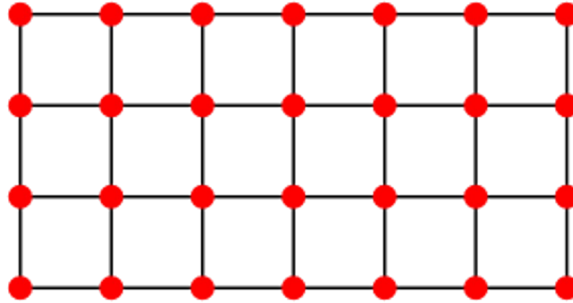
Cycle



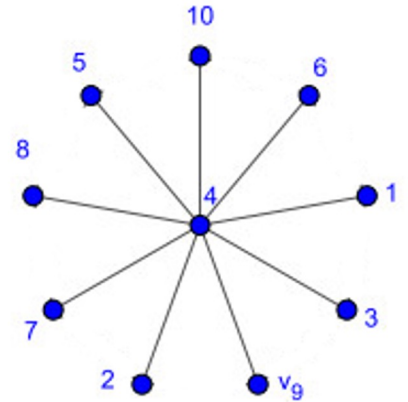
Line



Lattice



Star



These are not complex networks!

# Clique and Bi-partite clique

- A **clique** is a complete (sub)graph  $(E = (V \times V))$
- An **n-clique** is a complete graph of n nodes
- A **bi-partite clique** is such that

$$V = V_1 \cup V_2, V_1 \cap V_2 = \emptyset, E = (V_1 \times V_2)$$

- A **(n<sub>1</sub>, n<sub>2</sub>)-clique** is a bipartite clique such that

$$|V_1| = n_1, |V_2| = n_2$$

# Summary

# Things to remember

- . Definitions
  - degree, in-degree, out-degree, strength
  - time-evolving graph, multi-layer graph
  - line graph, cycle graph, star graph, lattice, bi-partite graph, clique
- . Writing the adjacency matrix of a graph, and drawing a graph given its adjacency matrix
- . Plotting the degree distribution of a graph
- . Projecting a bi-partite graph

# Sources

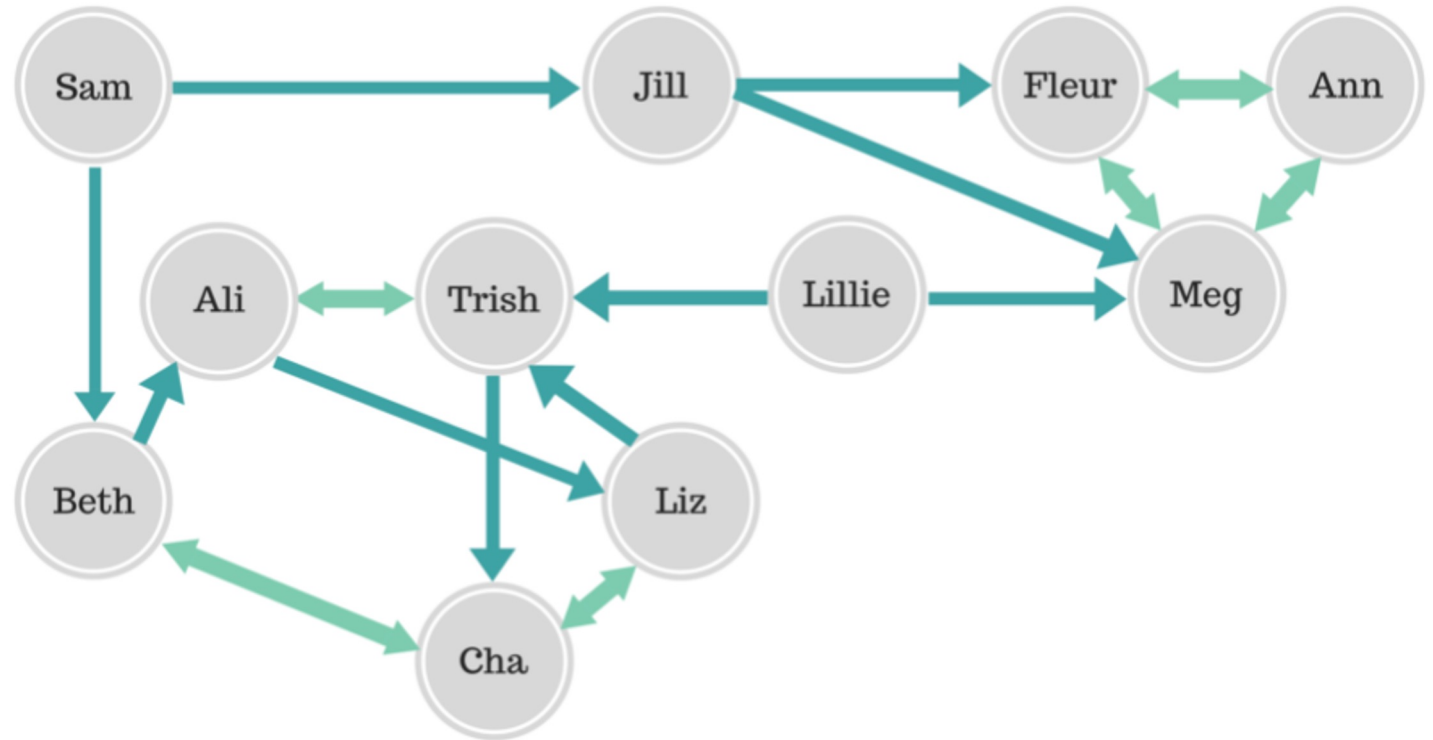
- A. L. Barabási (2016). Network Science – [Chapter 02](#)
- URLs cited in the footer of specific slides



# Practice on your own

Draw the  
indegree,  
outdegree,  
degree  
distribution

Write the  
adjacency  
matrix



# Practice on your own

How do you call the subgraph induced by nodesets:

- $\{H, A, B\}$
- $\{G, H, D\}$
- $\{B, D, E, G\}$
- $\{A, B, D, E\}$

